SCHEDULE INDEPENDENT SIMULATION: A NEW TECHNIQUE OF DISCRETE EVENT SIMULATION SUITABLE FOR SMALL COMPUTERS

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#### **ABSTRACT**

A state space approach to discrete event simulation is presented where an explicit mathematical model of the System is obtained using recursive mathematical functions similar to those utilized in continuous system simulation. Only the processes of stated interest in a simulation study are modeled.

Since events in the simulated system are not neccessarily executed in their real-time order of occurrence, many of the repeated list scans associated with conventional simulation techniques can be eliminated. This simulation algorithm constitutes an efficient simulation alternative that can be conveniently implemented in FORTRAN on today's smaller scale computers and time sharing systems.

#### I. BACKGROUND

Minicomputers and microcomputers constitute one of the fastest growing segments of the computer market (11). Therefore, attention focuses more keenly on development of simulation algorithms, such as the one described below, that are suitable for implementation on smaller scale computers.

Traditional simulation approaches such as GPSS, etc. contain built-in control algorithms to automatically monitor the sequencing of simulated events in their real-time order of occurrence. The resultant high storage and execution time requirements (particularly if the system is congested) generally preclude use of minicomputers.

A new technique of modeling discrete event systems, denoted SIS for Schedule Independent Simulation, was initially formulated at Northwestern University and constituted this author's recent Ph.D. thesis (8). By pinpointing simulation objectives, schedule independent techniques can be utilized. Minimum state space representations have been obtained for a variety of systems demonstrating serial and parallel connection. Computer implementations were in Fortran in batch mode on the CDC 6400 computer.

The development of SIS has been, and will continue to be, an ongoing project. Current work at the University of Toledo\* continues: (i) to extend the

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SIS algorithm to cover a broader class of queueing systems demonstrating feedback as well as serial and parallel connection; and (ii) to support development of related Fortran software. Recent computer implementations have been performed principally in an interactive mode on a 28K, PDP-11 time sharing system.

The following sections summarize the development of the SIS algorithm. It is in no way intended that SIS be viewed as constituting a "replacement" for conventional simulation analysis but simply as a convenient, efficient simulation "alternative" suitable to implementation on a small computer in a variety of "real world" queueing situations.

#### II. MODELING CONCEPTS

SIS is a simulation algorithm that:

- (i) models discrete event systems using explicit recursive mathematical functions (similar to those employed in continuous system simulation)
- (ii) processes events out of their real time order of occurrence and thereby eliminates expensive list processing associated with conventional time-sequential approaches
- (iii) models only the processes of stated interest in a particular simulation study, thereby achieving a minimum state space representation of the system
- (iv) can be conveniently implemented on a smaller scale computer or in the time sharing environment in a higher level language such as FORTRAN which is familiar to users and has attractive subprogram capabilities.

In many simulation studies the analyst can identify a very specific simulation objective, such as the estimation of a particular waiting time parameter or customer transit time, etc. SIS is particularly appropriate for modeling such systems, especially if a larger computer is not available or if the simulated system is congested. In some instances, execution time of the SIS recursive equations can be independent of the degree of congestion in the simulated system.

## MATHEMATICAL TECHNIQUES

A complex discrete event simulation model can be viewed as a queueing network with output from one queue comprising the input to another. Each queue is characterized by an arrival pattern, service mechanism, and a queue discipline. The queues in the model may be connected by serial, parallel, or feedback connection.

In a general state-space representation of a system S defined over T we have:

$$S \begin{cases} s(t) = f[s(to), \mu(to,t)] \\ y(t) = h[s(t), \mu(t)] \end{cases}$$
 (1)

where: to  $\underline{\Delta}$  starting point

 $s(t) \Delta state$  of the system at t  $\epsilon$  parameter set T.

 $\begin{array}{c} \mu(t) \stackrel{\Delta}{\underline{\triangle}} \text{ value of } \underbrace{\begin{array}{c} \text{input} \\ \text{input} \end{array}}_{f} \text{ function at } t \in T. \\ y(t) \stackrel{\Delta}{\underline{\triangle}} \text{ value of output function at } t \in T. \\ \text{input segment } (to,t) \stackrel{\Delta}{\underline{\triangle}} \text{ input up until } t \in T. \end{array}$ 

In a conventional GPSS or Simscript computer implementation, for example, the parameter set T represents time and processing throughout the simulation is time-sequential, accomplished via "scheduling" techniques. Lists of potential events are maintained, and as events are scheduled, they are placed on "chains" which are subsequentially scanned by the control algorithm and executed in their real-time order of occurrence.

The contribution of this work is to demonstrate how state representations can be achieved where the parameter set T does not represent time. Instead, times such as arrival times and departure times are viewed as "states" of the system and computed via recursive mathematical equations. Processing of some systems can be totally "schedule-independent", and execution time of the recursive mathematical equations can be independent of the degree of congestion in the simulated system. 

Example: As a simple example to demonstrate the modeling concepts involved, consider a system consisting of just one, single server, first-come-first-served (FCFS) queue. Let parameter set T  $\triangle$  {1,2,3,...} to represent arrival number into the system and not time. (Such processing shall be denoted as "customer-sequential processing"). Define input functions as:

- $\mathbf{u_1}(i) \triangleq interarrival time between i^{th}$  and  $(i+1)^{st}$  arrivals
- $u_2(i) \triangle length of service time of i<sup>th</sup> arrival$

Let states of the system be defined as:

TA(i)  $\underline{\underline{\Delta}}$  time of i<sup>th</sup> arrival into system TD(i)  $\underline{\underline{\Delta}}$  time of i<sup>th</sup> departure  $\underline{\underline{M}}(i)$   $\underline{\underline{\Delta}}$  waiting time in line of i<sup>th</sup> arrival th N(i)  $\underline{\underline{\Delta}}$  queue size upon and including the i arrival

Then, to study waiting times or queue size, we obtain the following recursive state transition functions:

$$TA(i+1) = TA(i) + u_1(i)$$
(3)

$$W(i+1) = MAX \{0; W(i) + u_2(i) - u_1(i)\}$$
 (4)

$$TD(i+1) = TA(i+1) + W(i+1) + u_2(i+1)$$
 (5)

$$N(i+1) = N(i) + 1 - \sum_{j=i+1-N(i)}^{i} B(j)$$
(6)

where 
$$B(j) = \begin{cases} 1 & \text{if } TA(i+1) \ge TD(j) \\ 0 & \text{otherwise} \end{cases}$$
  
 $TA(1)$  is given;  $N(1) = 1$ 

Processing in a SIS Simulation can be "customersequential", where a customer is processed through a major portion, if not the entire system, before processing of the next sequential customer to enter the system is even initiated (as demonstrated by the example). Also, as described in (8), processing can be "queue-sequential", which is attained by defining the parameter set T  $\triangle$  {1,2,3,...} to represent queue number in the system. Computation of the recursive state equations is performed by simulating all customers through any queue (i) before any customer is processed at queue (i+1).

# STATE SPACE MODELS FOR REPRESENTATIVE SYSTEMS

Systems depicting serial, parallel, and feedback connection have been analyzed and recursive state space representations achieved. Cases have included:

- single server, FCFS queues in tandem with both limited and unlimited line length, using both customer sequential processing and queue sequential processing.
- parallel server queues with unlimited line (ii) lengths
- (iii) priority, non-preemptive queue with one class priority customer
- (iv) mixture of above
- feedback queues with unlimited line lengths and unlimited number of feedbacks.

Earlier derivations and FORTRAN programs in (8) are confined to systems depicting only serial or parallel connection. Computer implementations were on a CDC 6400 computer. Recent work also emphasizes the modeling of feedback queues and the development of interactive FORTRAN programs executed on a 28K, PDP-11 time sharing system.

A few sample models are presented below to demonstrate the modeling theory.

#### 1. Serial Connection: FCFS Queues in Tandem

Many real world stochastic systems can be described by a series of queues in tandem depicting serial connections (or perhaps serial connection combined with some parallel connection). Moreover, it is in

the simulation of the serially connected, congested system that the SIS algorithm can make the most significant contribution.

It is here that an application of the SIS algorithm could result in totally schedule-free processing with all time consuming list scans eliminated. (See following example.) Execution time of the recursive state equations could be independent of the degree of congestion in the simulated system. A simulation model that might have significant storage requirements and execution time in a time-sequential algorithm such as GPSS could be reduced to a nominal FORTRAN problem.

<u>State Equations</u>: Consider a series of FCFS, single-server queues in tandem with unlimited line lengths. Let parameter set T  $\underline{\Delta}$  {1,2,3,...} to represent arrival number into system. Define state variables as:

$$TA(i)_{k+1} \triangleq time of arrival of customer (k+1) at queue (i)$$

$$W(i)_{k+1} \triangleq \text{waiting time of customer (k+1) at queue (i)}$$

$$\text{TD(i)}_{k+1} \triangleq \underset{\text{from queue (i)}}{\text{time of departure of customer (k+1)}}$$

$$N(i)_{k+1} \triangleq \begin{array}{l} \text{Number of customers at queue (i) upon} \\ \text{and including arrival of customer (k+1)} \\ \text{to the queue.} \end{array}$$

Input is:

$$S(i)_{k+1} \triangleq \text{service time of customer (k+1) at queue (i)}$$

Then the following state transition functions are obtained:

$$TA(i)_{k+1} = \begin{cases} TA(i)_k + I_k & \text{for } i=1 \\ TD(i-1)_{k+1} & \text{for } i>1 \end{cases}$$
 (7)

$$W(i)^{*}_{k+1} = \begin{cases} W(1)_{k} + S(1)_{k} - I_{k} & \text{for } i=1 \\ W(i)_{k} + S(i)_{k} - S(i-1)_{k+1} \\ + [\lambda_{i-1}, k+1] & W(i-1)^{*}_{k+1} & \text{for } i>1 \end{cases}$$
(8)

where 
$$\lambda_{i-1,k+1} = \begin{cases} 1 & \text{if } W(i-1)*_{k+1} < 0 \\ 0 & \text{otherwise} \end{cases}$$

$$W(i)_{k+1} = MAX \{0; W(i)^{*}_{k+1}\}$$
 (9)

$$TD(i)_{k+1} = TA(i)_{k+1} + W(i)_{k+1} + S(i)_{k+1}$$
 (10)

$$N(i)_{k+1} = N(i)_{k} + 1 - \sum_{j=k+1-N(i)_{k}}^{k} B(i)_{k+1}, j$$
 (11)

where: 
$$B(i)_{k+1,j} = \begin{cases} 1 & \text{if } TA(i)_{k+1} \ge TD(i)_{j} \\ 0 & \text{otherwise} \end{cases}$$

Computer Implementation: A conservative example consisting of five single server queues in tandem with unlimited line lengths was simulated in FORTRAN (for 3 levels of system congestion) using the SIS algorithm for customer-sequential processing. An associated GPSS simulation was also performed.

Results indicated that execution time of the SIS simulation remained <u>constant</u> (ie., independent of the degree of congestion in the simulated system). SIS simulations were about 5 times faster than the associated GPSS simulations even in this conservative example.

In our current environment, such a typical SIS simulation could be performed on a 28K, PDP-11 time sharing system while the GPSS simulation (submitted as a low priority job) requires use of the IBM 360 and a 108K partition.

#### 2. Parallel Connection

The SIS simulation algorithm has been applied to a multiple server queue. In this instance, some time-sequential processing is combined with the recursive techniques.

Consider a multiple server queue with "NS" servers. A customer is assigned to a specific server via an appropriate assignment algorithm. Let parameter set T represent arrivals. Define state variables as:

 $\mathsf{TA}_{k+1} \triangleq \mathsf{Time} \ \mathsf{of} \ \mathsf{arrival} \ \mathsf{of} \ \mathsf{customer} \ (k+1) \ \mathsf{at} \ \mathsf{the} \ \mathsf{queue}$ 

 $\mathsf{TD}_{k+1} \triangleq \mathsf{Time}$  of departure of customer (k+1) from the queue

 $W(i)_{k+1} \triangleq Waiting time of customer (k+1) in line (i)$ 

 $N(i)_{k+1} \triangleq \begin{array}{l} \text{Number of customers in line (i) and at server (i) upon and including the arrival of customer (k+1)} \end{array}$ 

Input is:

I(k)  $\underline{\triangle}$  Interarrival time between customer (k+1) and customer (k) at the queue

 $S(i)_{k+1} \triangleq Service time of customer (k+1) at server (i)$ 

Then the following state transition functions can be obtained:

$$TA_{k+1} = TA_k + I(k)$$
 (12)  
 $W(I)_{k+1}^* = \begin{cases} W(i)_{last} + S(i)_{last} - \sum_{j=last}^{k} I(j) \end{cases}$  (13)

$$W(I)*_{k+1} = \begin{cases} W(i)_{last} + S(i)_{last} - \sum_{j=last}^{\infty} I(j) \text{ (1)} \\ k \\ - \sum_{j=l}^{\infty} I(j) \text{ if customer (k+1) is the } \\ j=l & \text{first customer to use} \\ & \text{server (i)} \end{cases}$$

where "last" refers to the index number of the last previous customer to be assigned to server (i)

$$W(i)_{k+1} = MAX \{0; W(i)^*_{k+1}\}$$
 (14)

$$TD_{k+1} = TA_{k+1} + W(i)_{k+1} + S(i)_{k+1}$$
 (15)

$$N(i)_{k+1} = \begin{cases} N(i)_k + \lambda - NGONE(i, k+1) & (16) \\ where \lambda = \begin{cases} 1 & \text{if customer } (k+1) \\ assigned & \text{to line } (i) \end{cases}$$
or
$$\begin{cases} 1 & \text{if customer } (k+1) & \text{is the first } \\ customer & \text{to enter line } (i) \end{cases}$$

where: NGONE(i,k+1) △ Number of customers that departed server (i) after the arrival of customer (k) and before arrival of customer (k+1)

#### 3. Feedback Connection

Feedback connection has been more difficult to model using the SIS algorithm. A recursive model which also incorporates some time sequential techniques has been obtained for a single server FCFS queue in which customers have the option of leaving the queue upon completion of initial service or feeding back an unlimited number of times. Computer implementation is on a PDP-11 time sharing system where it is determined in an interactive mode whether a customer feeds back.

The capabilities of this model are being expanded and it will be presented in more detail at the conference.

#### III. COMPUTER IMPLEMENTATIONS

Since the major contribution of this work involves presenting a new discrete event modeling technique, emphasis in the computer examples has been on applying appropriate modeling concepts rather than on performing many simulations of real world systems and interpreting statistical results.

The iterative nature of the state equations makes computer implementations particularly well suited to the "horizontal" type of processing we have denoted as "customer sequential" processing. Therefore, customer sequential processing has been emphasized throughout the work.

To demonstrate the modeling theory, one typical simulation run for a series of FCFS queues in tandem is presented in the Appendix along with the associated Fortran source code. Other examples of simulations of queueing systems depicting both serial and parallel connection are presented in (8). Recent simulations of feedback queues (not described in (8)) have been performed on a PDP-11 time sharing system and are also available.

#### IV. FUTURE WORK PLANNED

Future work scheduled will extend the SIS algorithm to cover a broader class of queueing situations than those described in section II. Associated Fortran subprograms will be developed to offer the user a timely, efficient simulation alternative suitable to implementation on a smaller scale computer or in the time sharing environment.

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FORTRAN, IV VOIC-03
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PROGRAM GHOR(INPUT, OUTPUT)
               PROGRAM GHOR(IMPUT, GUTPUT)
PROGRAM GHOR SIMULATES A SERIES OF FCFS SINGLE SERVER QUEUES
IN TANDEM: LINE LENGTHS HAY BE INFINITE OR BOUNDED.
RECURSIVE STATE EQUATIONS ARE GIVEN FOR -- VAITING TIMES IN
QUEUES, NUMBERS IN QUEUES, ARRIVAL TIMES AT QUEUES, DEPARTURE
TIMES FROM QUEUES, STATE EQUATIONS ARE COMPUTED SEQUENTIALLY BY
CUSTOMER. THAT IS, COMPUTATIONS ARE PERFORMED HORIZONTALLY
ACROSS GRID, CALCULATING ALL QUEUE DATA FOR CUSTOMER(K) BEFORE
PROCESSING CUSTOMER(K+1) AT ANY QUEUE.
COMMON/BLI/MUMBER.KSTOP.MQ. ISIZE.TARRIU.ICAP(KQ)
   0001
                     COMMON/BLI/NUMBER, KSTOP, NQ, ISIZE, TARRIV, ICAP(40)
   0008
                     COMMON/BLE/NUMB( 20), W( 20), S( 20), TDEPRT, RINT, TASAVE
   0003
                     COMMON/ BL3/101,108
   0004
                     COMMON/BL4/VEC(20,20),L(20)
   0005
                     COMMON/PL5/SOLD(80)
   0006
                     COMMON/BL9/A, B, X, STDEV, XMU, Q, IX, IY
                COMMON/BLIO/IUNIT.JUNIT
IUNIT IS THE OUTPUT DEVICE USED.
   0007
                JUNIT IS THE INPUT DEVICE USED.
   0008
                    IUNIT=7
   0009
                     JUNIT=5
              WRITE(IUNIT, 600)
600 FORMAT(IH , 'INPUT DATA FOR THIS SIMULATION RUN AS
1 DEFINED. FOLLOW THE ALPHA')
WRITE(IUNIT, 602)
  0010
  1 100
  0018
  0013
              602 FORMATCIK , FORMAT SPECIFIED AND SEPARATE ALL VARIABLES
1 BY COMMAS. *)
              0014
  0015
   0016
                    WRITE(IUNIT, 606)
  0017
              606 FORMATCINO, "INPUT NRINS -- NUMBER OF RINS, FORMAT 110"
  0015
                    WRITE(IUNIT, 608)
  0019
              608 FORMATCIHO, " ")
  DORO
                    READCJUNIT, 100) NRUNS
  0051
              100 FORMATCIIO)
  0022
                    DO 89 ICOUNT=1,NRUNS
  0023
                    IOMO DE=1
              IOMODE-1
READ INPUT DATA FOR THIS SIMULATION RUN-
NUMBER --- RUN ID NUMBER
KSTOP -- NUMBER OF CUSTOMERS TO BE SIMULATED
NQ -- NUMBER OF QUEUES IN TANDEM
ISIZE -- EQUALS 1 IF LINES BOUNDED- O OTHER
TARRIV -- TIME OF ARRIVAL OF FIRST CUSTOMER
                                                                      O OTHERWISE
               TOUT -- OPTIONAL OUTPUT IS WRITTEN EVERY TOUT CUSTOMERS
               ICAP(I) -- CAPACITY OF LINE(I) IF LINES ARE BOUNDED
              FRITE(IUNIT, 612)
612 FORMAT(1HO, "NUMBER- RUN ID NUMBER, FORMAT IIU")
  0024
  0025
                    WRITE(IUNIT, 614)
  0026
              614 FORMAT(1H , KSTOP -- NUMBER OF CUSTOMERS TO BE SIMULATED,
1 FORMAT 110')
WRITE(1UNIT,616)
  U027
  -0025
              616 FORMAT(IN , "NQ-- NUMBER OF QUEUES IN TANDEM, FORMAT 110")
WRITE(IUNIT,618)
  0029
  0030
  UÓ31
              618 FORMATCH , "ISIZE -- EQUALS 1 IF BOUNDED, O OTHERWISE")
  0032
             WRITE(IUNIT, 620)
620 FORMAT(1H , "TARRIV-- TIME OF ARRIVAL OF FIRST CUSTOMER,
1 FORMAT F10.3")
  0033
 0034
                   WRITE(IUNIT, 622)
 0035
             622 FORMAT(IH , '10UT-- OPT- OUTPUT WRITTEN EVERY 10UT
1 CUSTOMERS, FORMAT 110')
 0036
             WRITE(IUNIT, 684)
684 FORMAT(IHO, 994)
 0037
                   READ(JUNIT, 101) NUMBER, KSTOP, NQ, ISIZE, TARRIV, IOUT
 0038
 0039
             101 FORMATC4110, F10.3, 110)
             IF(ISIZE.EQ.D) GO TO 799
WRITE(IUNIT,626)
626 FORMAT(IHO, INPUT CAPACITIES OF LINES SEPARATED BY COMMAS,
 0040
 0042
 0043
                 I FORMAT IIO"
             WRITE(IUNIT, 628)
626 FORMAT(INO, "7")
 0044
 0045
 0046
                  READ(JUNIT, 102)(ICAP(I), I=1,NQ)
 0047
             102 FORMATCC7110>>
              PARAMETERS FOR RANDOM NUMBER GENERATORS.
            799 WRITE(IUNIT, 640)
640 FORMAT(1HO, "INPUT PARAMETERS FOR RANDOM NUMBER
 0048
 0049
            0050
0051
U05#
0053
            WRITE(IUNIT, 646)
646 FORMAT(IHO, SERVICE TIME UNIFORMLY DISTRIBUTED OVER
0054
0055
                1 THE INTERVAL (A.B). 4)
            WRITE(IUNIT, 648)
648 FORMAT(IH , FORMAT F10.3")
0056
0057
            WRITE(IUNIT, 650)
650 FORMAT(IH, 'STDEV--STANDARD DEVIATION FOR NORMALLY
0058
0059
                I DISTRIBUTED RANDOM ")
0060
                  WRITE(IUNIT, 651)
            651 FORMATCIH ,
1800
                                              VARIABLE WHEN SERVICE TIME IS A
                1 FUNC. OF WAITING TIME, ")
WRITE(IUNIT, 652)
8900
0063
            658 FORMATCIH .
                                              FORMAT F1U.3"
0064
                 WRITE(IUNIT, 653)
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653 FORMATCIN . "X -- MEAN OF THE NORMALLY DISTRIBUTED
 0065
            1 RANDOM VARIABLE, FORMAT FID.3")
WRITE(IUNIT, 654)
654 FORMAT(IH, "XMU--MEAN OF THE NEGATIVE EXPONENTIAL
 0066
 0067
                 1 DISTRIBUTION FOR')
            WRITECIUVIT, 656)
0068
                                           INTERARRIVAL TIMES, FORMAT F10-3"
                  WRITE(IUNIT, 658)
 0070
            658 FORMAT(IH , 'Q--MINIMUM WAITING TIME BEFORE SERVICE
1 TIME BECOMES A FUNC. ')
 0071
             WRITE(IUNIT, 660)
 0072
 8073
                                        OF THE WAITING TIME AT THE PREVIOUS
            1 QUEUE, FORMAT F10.3°)
WRITECUNIT,661)
'661 FORMAT(1HO,°°)
 0074
 0075
 0076
                   READCJUNIT, 668) A, B, STDEV, X, XMU, Q
 0077
             662 FORMAT(6F10-3)
 0078
                  IX=0
              17=0
DO 10 I=1,100
10 CALL RANDU(IX,IY,YFL)
 UQ.79
 0050
 0061
              INITIALIZATIONS
             SCI) IS SERVICE TIME AT QUEUE(I). TDEPRT IS DEPARTURE TIME
FROM QUEUE(I). W(I) IS WAITING TIME AT QUEUE(I). NUMB(I) IS
NUMBER OF CUSTOMERS IN QUEUE(I) UPON AND INCLUDING ARRIVAL OF
CURRENT CUSTOMER.
 008 ¥
                  CALL OUTPTCIOMODE, K. I. IGUT)
0083
0064
                  IOMODE=2
IO1=NQ*(IOUT-1) +1
IO2=NQ*(IOUT)
 0085
 0086
                   TASAVE=TARRIV
                  DO 8 1=1.NO
 0087
 0088
 0089
                  NUMB(I)=1
 0090
               2 V(I)=0.
 U09 1
                  K=1
 0092
                  CALL OUTPT(IONODE)
 11003
                  CALL ANALYSCIOMODE, K, IOUT)
 0094
                  IOMODE=3
              COMPUTE STATE EQUATIONS FOR CUSTOMER(1) AT ALL QUEUES
                  DO 3 I=1,NQ
IF(I.EQ.1)GO TO 9
0095
 UD96
 0098
                  S(I)=SGET(I,W(I-1))
 0099
                  GO TO 13
S(I)=SGET(I)
 0100
 0101
              13 TDEPRT=TARRIV+V(1)+S(1)
0105
                  IPOS=L(I)+1
                  CALL PUSHUP(2, TDEPRT, 1, 1POS, NOFF) ... CALL OUTPT(10MODE, K, 1)
0103
0104
0105
                  TARRIV=TDEPRT
             CALL ANALYSCIONODE, K, IOUT)
MAJOR LOOP -- COMPUTE STATE EQUATS. FOR CUST. (K), FOR K .GT. 1
0106
0107
                  DO 44 K=2,KSTOP
               DO 5 I=1,NQ
5 SOLD(I)=S(I)
ULOR
0109
         C
              GET INTERRARIVAL TIME RINT
0110
                  RINT=RGET(K)
              INTERIOR LOOP-- COMPUTE STATE EQUATIONS FOR CUST-(K)
              AT EACH QUEUE.
0111
                  DO 4 I=1,NQ
TARRIV=TIMEA(1)
0112
             IF (ISIZE.EG.1) GO TO 6
CHECK FOR PREVIOUS DEPARTURES SO , OF CUSTOMERS CAN BE COMPUTE
7 CALL PUSHUP(1,TARRIV.1,1POS,NOFF)
0113
0115
0116
                  NUMB(I)=NUMB(I)+1-NOFF
0117
               5 IF(I.EQ.1) TASAVE=TARRIV
              GET WAITING TIME OF CURRENT CUSTOMER IN QUEUE(1)
0119
                  W(I)=WATE(I)
             IF(1.EQ.1)GO TO 12

GET SERVICE TIME OF CURRENT CUSTOMER AT QUEUE(1)

S(1)=SGET(1,W(1-1))
0120
         C
0192
             12 S(1) = SGET(1)
COMPUTE DEPARTURE TIME OF CURRENT CUSTOMER FROM QUEUE(1)
0124
       · C
0125
              14 TDEPRT=TARRIV+W(1)+S(1)
             IPOS=L(1)+1
ADJUST TD-VECTOR
CALL PUSHUP(2, TDEPRT, T, IPOS, NOFF)
CALL OUTPT(IOMODE, K, I)
0126
0127
0128
                  GO TO 4
0129
               TIME LENGTHS NOT INFINITE. CHECK TO SEE IF BOUNDING OCCURS.
6 CALL PUSHUP(3,TARRIV,1,1POS,NOFF)
SURPLS=NUMB(1)-NOFF-1CAP(1)
              LINE LENGTHS NOT INFINITE.
0130
0131
0138
                  IF(SURPLS-LE-0-) GO TO 7
             BOUNDING OCCURS

ADJUST SERVICE TIME OF CURRENT CUSTOMER AT PREVIOUS SERVER OR, IF I=1, DELAY ARRIVAL TIME OF CURRENT CUSTOMER.
         C
            DUMMY-BLOCK(I)
WRITE(IUNIT, 700)K, I
700 FORMAT(IH , "CUSTOMER", I4, "BLOCKED AT QUEUE", I4)
0134
0135
0136
            IF (I-NE-I) WRITECIUNIT, TOWN DUMMY

70% FORMAT(IK , "SERV. TIME & DEPRT. TIME AT PREVIOUS QUEUE

1 ARE INCRMTED. BY ", F6. %)

GO TO 7

4 CONTINUE

44 CALL ANALYSCIOMODE, K, IOUT)
0137
0139
0140
0141
0148
                  IOMODE=4
CALL OUTPT(IOMODE)
0143
0144
0145
             89 CALL ANALYSCIOMODE)
U146
                  STOP
0147
                  END
```

IOUT CUSTOMERS 3 IO=IO+1

10=0

RETURN

RETURN

END

IFCIO-LT-101) RETURN

IF(10-NE-102) RETURN

105 FORMATC 11H END OF RUN)

4 WRITECIUNIT, 105)

WRITECIUNIT, 104) N. I. TARRIV, NUMB(I), W(I), S(I), TDEPRT 104 FORMAT(8110, F10.3, 110, 3F10.3)

0033

0036 0037 0038

0040

0041

0042

0043

0044

0045

```
SUBROUTINE PUSHUP(MODE, T, I, IPOS, NOFF)
THIS IS THE LIST PROCESSOR WRITTEN TO HANDLE UPDATING OF TD-
VECTORS. TD-VECTORS ARE VECTORS OF SCHEDULED DEPARTURE TIMES
FOR EACH QUEUE. L(I) IS LENGTH OF ITH HOW OF MATRIX VEC.
              IN ALL MODES NOFF = NUMBER OF ELEMENTS TO BE PUSHED OFF TOP OF VECTOR VEC(1).
              MODE=1 --- COMPARE T WITH ROW I OF VEC. REMOVE ELEMENTS
MODE=2 --- ADD T TO VEC(I) IN POSITION (IPOS) IN VECTOR.
MODE=3 --- OBTAIN PUSHUP VALUE ONLY.
                                                                           REMOVE ELEMENTS .LE. T
                  COMMON/BL4/VEC( 20, 20), L( 20)
 0002
 0003
                  LL=L(I)
 0004
                  GO TO (1,2,1), MODE
 0005
               1 IF(LL.GT.0) GO TO 85
                  NOFF-0
 0007
 0008
                  RETURN
              25 DO 55 J=1,LL
IF (T-LT-VEC(1,J)) GO TO 66
 0009
 0010
 0012
              55 CONTINUE
 0013
                  J=J+1
 0014
              66 NOFF-J-1
 0015
                  IF(NODE-EQ-3) RETURN
                  IF(NOFF.EQ.D) RETURN
L(I)=L(I)-NOFF
IF(L(I)-GT.O) GO TO 17
 0017
 0019
 0080
 0022
                  RETURN
 0023
               2 L(I)=L(I)+1
 0024
                  IF(IPOS-GT-(L(I)-1)) GO TO 90
 0026
                  LL=L(I)
 0027
                  JJ=1 POS+1
             DO 91 J=LL,JJ
91 VEC(I,J)=VEC(I,J-i)
 98800
 0029
 0030
                  VEC(I, IPOS)=T
            90
0031
                  RETURN
 0032
             17 LL=L(I)
               DO 7 J=1,LL
7 VEG(1,J)=VEG(1,J+NOFF)
 0033
 0034
 0035
                  RETURN
 0036
                  END
FORTRAN IV
                        V01C-03
0001
                  SUBBOUTINE OUTPT(MODE, K, I, IOUT)
            OUTPUT WRITER (PROGRAM WRITES OUTPUT EVERY IOUT CUSTOMERS)
COMMON/BLI/NUMBER, KSTOP, NQ. ISIZE, TARRIV, ICAP(20)
0002
                  COMMON/BLE/NUMB(80), V(80), S(80), TDEPRT, RINT, TASAVE
0003
0004
                  COMMON/BL3/101,108
0005
                  COMMON/BLIO/IUNIT, JUNIT
0006
            GO TO (1,2,3,4),MODE
WRITE INITIAL CONDITIONS
1 WRITE(IUNIT,108)
0007
0008
            108 FORMAT(////)
0009
                  WRITECIUNIT, 100)NUMBER
           100 FORMAT(1H, "RUN NUMBER" ",13)

WRITE(IUNIT,500) NQ

500 FORMAT(1H,13," SINGLE SERVER FCFS QUEUES IN TANDEM")

WRITE(IUNIT,501) TARRIV

501 FORMAT(1H, "TIME OF ARRIVAL OF CUSTOMER(1)=",F5.2)
0010
0011
0012
0013
0014
0015
                 WRITE(IUNIT, 502) KSTOP
            502 FORMATCIH . 14. " CUSTOMERS SIMULATED")
·0016
           WRITE(IUNIT, 503) IOUT
503 FORMAT(1H , *OUTPUT WRITTEN EVERY *, 14, * CUSTOMERS *)
WRITE(IUNIT, 504)
0017
0015
00,18
            504 FORMATCH , "HORIZONTAL SINULATION SEQUENTIAL BY CUSTOMER")
0020
0021
                 IF(ISIZE, EQ. 1) GO TO ME
0023
                 WRITECIUNIT, 101>
0024
           101 FORMAT(82H INFINITE LINE LENGTHS)
0025
                 RETURN
0086
             88 WRITE(IUNIT, 102)(ICAP(II), II=1,NQ)
0027
            102 FORMATCH , "BOUNDED LINES" / BOUNDS ARE=",(7110))
0025
                 RETURN
           VRITE HEADINGS
2 WRITE(IUNIT, 103)
103 FORMAT(/5X,5H CUST, 4X, 6H QUEUE, 3X,7H TARRIV, 3X,7H NUMBER, 6X
12H W,8X,2H S,5X,7H TDEPRT)
0029
0030
0031
                 10-101-1
0032
                 RETURN
            WRITE OPTIONAL OUTPUT DATA FOR CUSTOMER(K) AT QUEUE(I) EVERY
```

```
0001
                 FUNCTION BLOCK(I)
             CURRENT CUSTOMER IS BLOCKED AT QUEUE(I). ADJUST S(I-1) OR RINT COMMON/BL1/NUMBER, KSTOP, NQ, ISIZE, TARRIV, ICAP(2U)
 0002
 0003
                 COMMON/BLE/NUMB(20), W(20), S(20), TDEPRT, RINT, TASAVE
             COMMON/BL4/VEG(20,20),L(20)
BLOCK = LENGTH OF TIME CUSTOMER IS BLOCKED AT QUEUE(I)
BLOCK=VEC(I,1)-TARRIV
 0004
 0005
                 IF(I.EQ.1) GO TO 1
 0006
             CUSTOMER MAINTAINS USE OF PREVIOUS SERVER WHILE BLOCKED AT
             QUEUE( I.) .
 0008
                 S(1-1)=S(1-1) + BLOCK
                 TDEPRT=TDEPRT+BLOCK
 0009
                 LL=L(I-1)
VEC(I-1,LL)=TDEPRT
 0010
 0011
                 TARRIV=TARRIV+BLOCK
 0012
  0013
                 RETURN
             CUSTOMER BLOCKED AT QUEUE(1) -- DELAY ARRIVAL INTO SYSTEM
 0014
               1 RINT=RINT+BLOCK
TARRIV=TIMEA(I)
 0015
 0016
                 RETURN
 0017
                 END
  0001
                 FUNCTION TIMEACL)
             PROGRAM COMPUTES TIME OF ARRIVAL OF CURRENT CUSTOMER AT
             QUEUE(1).
COMMON/BLB/NUMB(80), W(80), S(80), TDEPRT, RINT, TASAVE
  0002
                 IF(I.GT.1) GO TO 2
TIMEA=TASAVE+RINT
  0003
  0005
  0006
                 RETURN
  0007
               2 TIMEA-TDEPRT
 0008
                 RETURN
 0009
                 END
- 0001
                FUNCTION SGET(I,W)
            PROGRAM COMPUTES SERVICE TIME AT CURRENT QUEUE.
             THIS PROGRAM SHOULD BE TAILORED TO MEET REQUIREMENTS OF EACH
            SIMULATION.
COMMON/BL9/A, B, X, STDEV, XMU, Q, IX, IY
CALL RANDU(IX, IY, YFL)
 0002
 0003
 0004
                 IF(I.EQ. 1) 90 TO 1
            GO TO 2
SERVICE TIME AT QUEUE(1) IS A UNIFORMLY DISTRIBUTED RANDOM
 0006
         C
             VARIABLE ON INTERJAL (A.B)
              1 SGET=A+YFL*(B-A)
 0007
                RETURN
 0008
             SERVICE TIME IS A FUNCTION OF WAITING TIME AT PREVIOUS QUEUE #
             THE SERIES.
 0009
              2 SGET=2
                IF(W.LT.Q) GO TO 3 RETURN
 0010
 0012
             SERVICE TIME AT QUEUE(I) IS A NORMALLY DISTRIBUTED RANDOM VARIABLE WITH STANDARD DEVIATION "STDEY" AND MEAN "X".
 0013
              3 CALL GAUSS(ANORML)
                SGET=ANORML
 0014
 0015
                RETURN
 0016
                END
 0001
                 SUBROUTINE GAUSS(ANORML)
             SUBROUTINE GAUSS IS A DIRECT METHOD OF DERIVING NORMALLY
DISTRIBUTED RANDOM NUMBERS, BASED UPON THE PROPERTIES OF
RANDOM NUMBERS RATHER THAN AN EVALUATION OF THE INVERSE
             FUNCTION.
 0002
                 COMMON/BL9/A, B, X, STDEV, XMU, Q, 1X, 1Y
 0003
                 P=0.0
 0004
                 DO 10 IR=1,18
CALL RANDU(IX,IY,YFL)
P=P+YFL
 0005
 0006
 0007
             10 CONTINUE
 8000
                 ANORML=(P-6.0) *STDEV+X
 0009
                 RETURN
```

0010

```
0001
                  FUNCTION RGET(K)
             PROGRAM COMPUTES INTERARRIVAL TIME BETWEEN CURRENT AND PREVIOU
         C
                 CUSTOMER.
                 ARRIVALS ARE RANDOM
         C INTERARRIVAL TIMES GENERATED FROM NEGATIVE EXPON. DISTRIBUTION.
COMMON/BL9/A, B, X, STDEV, XMU, Q, IX, IY
CALL RANDU(IX, IY, YFL)
RGET=(-XMU*ALOG(YFL))
 0002
 0003
 0004
 0005
                  RETURN
 0006
                  END
             FUNCTION WATE(I)
PROGRAM COMPUTES WAITING TIME W(I) OF CURRENT CUSTOMER AT
QUEUE(I). W(I) DOES NOT INCLUDE SERVICE TIME.
 0001
 0002
                  COMMON/BLE/NUMB(80), W(20), S(20), TDEPRT, RINT, TASAVE
 0003
                  COMMON/BL5/SOLD( 80)
 0004
                  REAL LAMDA
 0005
                  LAMDA=0.
 8000
                  IF(1.GT.1) GO TO 6
                  WSTAR=W(1)+SOLD(1)-RINT
 0006
               GO TO 7
6 IF(WSTAR-LT-0-) LAMDA=1-
 0009
 0010
 0012
                  WSTAR=W(I)+SOLD(I)-S(I-1)+LAMDA+WSTAR
 0013
               7 W(I)=0.
 0014
                  IF(WSTAR. GT. 0.) W(1)=WSTAR
                  WATE=W(I)
 0016
 0017
                  RETURN
 0018
                  EN D
            SUBROUTINE ANALYS(MODE,K, IOUT)
PROGRAM COMPUTES CUSTOMER STATISTICS -- TRANSIT TIME.
AVERAGE CUSTOMER TRANSIT TIME IS COMPUTED.
0001
         C
                  COMMON/BLI/NUMBER, KSTOP, NQ, ISIZE, TARRIV, ICAP(80)
0002
                  COMMON/BLE/NUMB(80), W(80), S(20), TDEPRT, RINT, TASAVE
0003
0004
                  COMMON/BLIO/IUNIT, JUNIT
             TRANS-TRANSIT TIME FOR CURRENT CUSTOMER. TSUM-CUMULATIVE
         C
         C
             TRANSIT TIME.
              GO TO (2,2,3,4),MODE
3 TRANS=0.
0005
0006
0007
                  DO 1 I-I-NO
               1 TRANS=TRANS+S(1)+V(1)
TSUM=TSUM+TRANS
0008
0009
                  10=10+1
 0010
                  IF(10.NE-10UT) RETURN
 0011
 0013
                  10=0
 0014
                  WRITE(IUNIT, 100)K, TRANS
            100 FORMATCH , TRANSIT TIME FOR CUST', 13, " =", F10.3/)
 0015
0016
                 RETURN
               2 IO=10UT-1
 0018
                  TSUM=0.
 0019
                  RETURN
 0080
               4 AVT=TSUM/KSTOP
            WRITE (UNIT, 101) AUT
101 FORMAT(1H , 'AUGE TRANSIT TIME= ',F10-3)
 0081
 0022
 0023
                  RETURN
 0024
RUNSEXEC
 *QHORFL/CORE: 25
INPUT DATA FOR THIS SIMULATION RUN AS DEFINED. FOLLOW THE ALPHA FORMAT SPECIFIED AND SEPARATE ALL VARIABLES BY COMMAS.
INPUT NRUNS--NUMBER OF RUNS, FORMAT I10
1
NUMBER-- RUN ID NUMBER, FORMAT I10
KSTOP-- NUMBER OF CUSTOMERS TO BE SIMULATED, FORMAT I10
NQ-- NUMBER OF QUEUES IN TANDEM, FORMAT I10
ISIZE-- EQUALS I IF BOUNDED, 0 OTHERWISE
TARRIV-- TIME OF ARRIVAL OF FIRST CUSTOMER, FORMAT F10.3
IOUT-- OPT- OUTPUT WRITTEN EVERY IOUT CUSTOMERS, FORMAT I10
201,101,5,1,3,,20
INPUT CAPACITIES OF LINES SEPARATED BY COMMAS, FORMAT 110
4, 2, 2, 1, 2
INPUT PARAMETERS FOR RANDOM NUMBER GENERATORS AS DEFINED. FOLLOW
```

THE ALPHA FORMAT SPECIFIED AND SEPARATE ALL VARIABLES BY COMMAS.

Winter Simulation Conference

SERVICE TIME UNIFORMLY DISTRIBUTED OVER THE INTERVAL (A,B),

STORMAT F10-3
STDEV--STANDARD DEVIATION FOR NORMALLY DISTRIBUTED RANDOM
VARIABLE WHEN SERVICE TIME IS A FUNC. OF WAITING TIME. FORMAT F10-3

Y-MEAN OF THE NORMALLY DISTRIBUTED RANDOM VARIABLE, FORMAT F10-3
X--MEAN OF THE NEGATIVE EXPONENTIAL DISTRIBUTION FOR
INTERARRIVAL TIMES, FORMAT F10-3
Q--MINIMUM WAITING TIME BEFORE SERVICE TIME BECOMES A FUNCOF THE WAITING TIME AT THE PREVIOUS QUEUE, FORMAT F10-3

20,40,10,30,40,105

RUN NUMBER#201 MAN NUMBER=201
5 SINGLE SERVER FCFS QUEUES IN TANDEM
TIME OF ARRIVAL OF CUSTOMER(1) = 3.00
101 CUSTOMERS SIMULATED
OUTPUT WRITTEN EVERY 20 CUSTOMERS
HORIZONTAL SIMULATION SEQUENTIAL BY CUSTOMER
BOUNDED LINES BOUNDS ARE= 2

			******	¥	s	TDEPRT
CUST	QUEUE	TARRIV	NUMBER	-	_	5.986
1	1	3-000	1	0.000	2-986	8 686
1	.2	5-986	1	0.000	8- 699	13.670
1	3	8 686	.1	0-000	4.984	
1	4	13-670	1	0.000	2-845	16-514
1	5	16-514	1	0.000	2.544	19-058
TRANSIT TIME	FOR CUST	1 =	16-058			
	TH OOVER A	T AIRLIE	1			
	BLOCKED A		4			
		I QUEUE IME AT PI		ADE INC	RMTED. BY	1.05
SERV. TIME &			3	4.875	3. 731	81.990
21	1	73-385	ĭ	0.000	2.000	83.990
81	2	81.990	8	1.543	3.166	85 • 699
21	3	83-990	0	0.000	2.000	90.699
21	4	88 • 699	_	1.016	1.338	93.053
21	5	90.699	8	1.010	1.230	73-033
TRANSIT TIME	FOR CUST	51 =	19-668			
41	1	165.976	0	0.000	3 • 027	169 - 003
41	ģ	169-003	ĭ	0.000	3.417	172-420
41	3	178-420	ò	0.000	0.922	173.341
41	Ä	173-341	- ŏ	0.000	3-194	176-535
41	5	176-535	ŏ	0.000	2.005	178 - 540
TRANSIT TIME			12-564	0-00-		
INWASTI	, FUR COST	<b></b>	161 204			
61	1	230-590	3	5.750	8.868	238 • 602
61	2	238 - 602	ĭ	0.000	2.000	240-602
61	3	240.602	2	0.779	2.838	244-220
61	ă	244. 220	2	1-500	4.375	250-094
61	<u> </u>	250-094	ī	0.683	3.310	254-027
TRANSIT TIME			23-437			
1196404						
CUSTOMER 64	BLOCKED A	T QUEUE	1		·	
81	1	305- 206	0	0.000	3.351	306 - 557
81	2	308 • 557	1	0.000	3 • 648	312-205
51	3	312-205	0	0•.000	2-424	314-630
81	4	314-630	Q	G• 000	3- 737	318.366
81	5	315.366	Ö	0.000	2.249	320-615
TRANSIT TIM	FOR CUST	51 =	15-409			
		•	_			001 455
101	1	381-410	3	6.506	3. 257	391-473
101	2	391-473	0	0.000	2.000	393-473
101	3	393-473		0.136	2.407	396-015
101	4	396-015	8	2.535	2.549	401.399

404-335

8.000

0.936

END OF RUN AVGE TRANSIT TIME-STOP --19.330

TRANSIT TIME FOR CUSTIO! -

101

401 - 399

22.925

READY

BYEF