# SIMULATION OF LARGE WIRELESS SENSOR NETWORKS USING CELL-DEVS

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## ABSTRACT

The advancement of electronic sensing devices, microcomputers and wireless communication devices has lead to creation of new smart sensors, which can monitor actuate, compute and communicate. Typically, these sensors are deployed in non-deterministic mode (randomly) when deployed in large numbers. These sensor devices have the capability to self-organize into the so-called Wireless Sensor Networks (WSN). WSN are ad-hoc networks, consisting of these spatially distributed sensing and processing devices. We introduce a model and a simulation study of these Large Wireless Sensor Networks (WSN) by implementing the Topology Control Algorithm. We use the Cell-DEVS formalism, which enables efficient execution of cellular models. Thereafter, we observe and evaluate the behavior of sensor nodes and entire WSN from the simulation results obtained, under different test scenarios.

## 1    INTRODUCTION

The emergence of powerful embedded micro-computer systems for wireless sensor networks provides a good ground for creation of new smart sensor systems, which can be useful to further promote new scientific endeavors and enhance our lives.

The wireless sensor devices can monitor actuate, compute and communicate, yet are small in size and low in cost. The WSN are ad-hoc networks, consisting of these spatially distributed sensing and processing devices (Cunha, Silva, and Loreiro, 2005; Healy, Newe, and Lewis 2008). WSN are used in many different applications, such as medicine, transportation and urban monitoring, traffic control, military, environment and habitat monitoring, energy management, smart homes, industrial applications, etc. The effectiveness of WSN is not just in their monitoring, actuating, computing and communications capabilities: with the added processing power, analog and digital ports, transceivers and memory, they can self-organize and communicate in the deployed area (as depicted in Figure 1). Their processing power is limited however, and WSNs are usually deployed in large numbers and their load is shared accordingly.
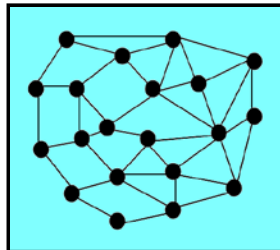


Figure 1: Sensor nodes self-organized

Due to the fact that sensor nodes have limited bandwidth, computing power and limited energy resources, one of the constraints in WSN is the energy efficiency of sensor nodes (i.e., their power consumption). There are many different approaches to solving the energy efficiency problem, such as the Topology Control Algorithm (Cunha, Silva, and Loreiro 2005; Shen, and Guo 2007; Ye et al. 2002). The objective of this algorithm is related to the efficiency of WSN network, and it fo-

cused on how to increase its lifetime. The rationale of the network topology control algorithm is to reduce the number of redundant sensor nodes monitoring a particular region, hence increasing the efficiency and lifetime of the WSN network. In a nutshell, topology control exploits the redundant deployment of the sensor nodes, overcoming their energy limitations by restricting the set of nodes which are considered neighbors of a given node, while making sure that sensing area is still covered by a sufficient number of sensors. Furthermore, it reduces interference problems (which are noticeable when large number of sensor nodes are active).

We here introduce a model and a simulation of Wireless Sensor Networks (WSN) by implementing the Topology Control Algorithm using the Cell-DEVS formalism (which enabled efficient execution of this cellular model). The complexity of the problem can be simplified using the CD++ toolkit and the Cell-DEVS formalism (implementing this simulation in a other high level programming language is much more complex).

The rest of the paper is organized as follows: we first give an overview of the Cell-DEVS formalism, followed by the WSN model definition. Thereafter, we observe and evaluate the behavior of sensor nodes and entire WSN from the simulation results obtained, under different test scenarios. We then describe and discuss the simulation results, analysis and discussion of the initial and improved results for different scenarios.

## 2    MODEL DEFINITION

Due to the complexity of the model under study, we used Cell-DEVS (Wainer 2009) and the CD++ toolkit (Wainer 2002), as an efficient way to model and simulate cellular models (Chopard and Droz1998), in our case the WSN topology problem. Cell-DEVS is an extension to the DEVS formalism (Zeigler, Kim, and Praehofer 2000), which has been used to model systems that can be represented as cell spaces. A Cell-DEVS model is represented as a cell space, where each cell is represented as an atomic DEVS model. Each cell is connected to the local neighboring cells. A delay mechanism in each cell (transport delay or inertial delay) is used to delay the propagation of state change events through the cell space, providing the means for defining complex temporal behavior. An Atomic Cell-DEVS can be defined as follows:

$$TDC = < X, Y, I, S, \theta, N, d, \delta int, \delta ext, \tau, \lambda, D >$$

Where X is the set of external input events; Y is the set of external output events; I represents the definition of the model's modular interface; S is the set of possible states for a given cell; $\theta$ is the definition of the cell's state variables, N is the set of values for the input events; d is the delay of the cell; $\delta int$ is the internal transition function; $\delta ext$ is the external transition function; $\tau$ is the local computing function; $\lambda$ is the output function, and D is the duration function.

A Coupled Cell-DEVS model is built by connecting a number of Atomic Cell-DEVS models together into a cell space of any shape (including 2D and 3D cell spaces). The borders of the cell space can be either wrapped, in which case the cells at the border from one side of the cell space are considered neighbors to the cells at the border on the opposite side of the cell-space, or non-wrapped, in which case the border cells must have special rules defined by the modeler. A formal definition of Coupled Cell-DEVS is:

$$GCC = < Xlist, Ylist, I, X, Y, n, \{t1,….,tn\}, N, C, B, Z, select >$$

Where Xlist is the input coupling list; Ylist is the output coupling list; I represents the definition of the model's modular interface; X is the set of external input events; Y is the set of external output events; n is the dimension of the cell space; {t1,….,tn} is the number of cells in each of the dimensions; N is the neighborhood set; C is the cell space; B is the set of border cells; Z is the translation function; and select is the tie-breaking function;

CD++ is an M&S toolkit that implements DEVS and Cell-DEVS theory ((Wainer 2009; Wainer 2002; Zeigler, Kim, and Praehofer 2000). Atomic models can be defined using a state-based approach (coded in C++ or an interpreted graphical notation), while coupled and Cell-DEVS models are defined using a built-in specification language. CD++ also includes an interpreter for Cell-DEVS models. The model specification includes the definition of the size and dimension of the cell space, the shape of the neighborhood and borders. The cell's local computing function is defined using a set of rules with the form: POSTCONDITION DELAY {PRECONDITION}. This indicates that when the *PRECONDITION* is satisfied, the state of the cell will change to the designated *POSTCONDITION*, whose computed value will be transmitted to other components after consuming the *DELAY*.

We have used these facilities to create an advanced model of WSN, in which we can analyze the WSN topology problem. This problem is closely related to the minimum configuration of nodes for fully operational WSN, taking into account sensor node energy limitations for a long lasting - survivable WSN networks. In our model, we have structured the area as a two-dimensional cell space of size *n* x *n*, where every cell represents one sensor node. The WSN considered is such as all the sensor nodes have same properties (homogeneous WSN) and flat (i.e. no hierarchy among nodes). Each node can have up to 8 neighbors and there are *3 possible states* for each cell (*active*, *stand-by*, and *inactive*). For this model, Moore's neighbor-

hood is adopted (i.e., the origin cell and its 8 close neighbors). As a result, the Cell-DEVS simulation model, gives an insight into an elegant way of implementation the Topology Control Algorithm for a large WSN; in our particular case, addressing the issue of sensor node energy conservation to achieve a longer lifetime operation of the WSN.

During the *active mode* of operation, the node emulates an active sensor within the WSN (i.e. performs processing tasks, hence using energy which decreases with time). In the beginning of simulation, all the sensor nodes deployed, have the same amount of energy which decreases as time progresses while node is in active mode (maximum energy consumption) or in stand-by mode (minimum energy consumption). During the *stand-by mode*, sensor node is consuming a minimal amount of energy; it wakes up randomly in order to see if other close by nodes are already sensing/monitoring the predefined neighborhood area. If less than two sensor nodes within the neighborhood are *active*, the cell goes again into *stand-by* mode; otherwise it becomes *active*. After the entire energy of a cell is consumed, the cell becomes *inactive*; this process continues until all the sensor nodes' energy is consumed (i.e. all the cells are inactive).

The **n x n** cells in the cell space considered were organized in two planes reflecting a three dimensional space implemented to meet the basic constraints (of the defined WSN Topology problem) while exploring and capturing the main tasks of the problem considered. Each sensor node starts to operate with a fixed amount of energy. In this model the energy levels adopted based on (Cunha, Silva, and Loreiro 2005; Ye et al. 2002) are the following:

- WSN sensor energy (at the beginning of operation) = 0.8 J
- The energy of an WSN sensor node in *active mode* decreases by 0.0165 J every time step (in our case every 1 sec)
- The energy of a WSN sensor node in *stand-by mode* decreases by 0.00006 J every time step (in our case every 1 sec)

The WSN sensor node possible states are:

**Plane 0:**

**2** - WSN sensor node is in *active mode* within the neighborhood (WSN sensor coverage area)
**1** - WSN sensor node is in *stand-by mode* within the neighborhood (WSN sensor coverage area)
**0** – WSN sensor node is *passive* (i.e. energy of a node is consumed)

**Plane 1:**

**0.8** - initial energy level of WSN sensor nodes
**-1** - WSN sensor node is *passive* (i.e. energy of a node is consumed)
Neighborhood ={ (-1,-1,0) (-1,0,0), (-1,1,0), (0,-1,0), (0,0,0), (0,1,0)( 1,-1,0),(1,0,0),(1,1,0), (0,0,1), (0,0,-1) }

Plane 0 contains the different deployed sensor nodes, and it is used to be observed throughout the simulation of the model. Plane 1 was used as "memory" for keeping track of the energy levels of active and stand-by nodes. The correct decrease of energy level and node behavior in the Plane 0 is interrelated to Plane 1 which contains the sensors energy information throughout the simulation.

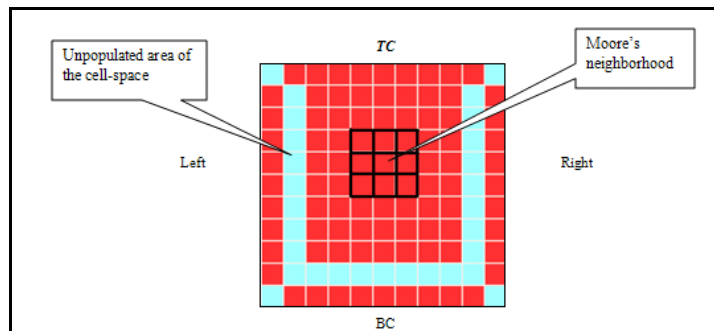The Plane 0 areas are organized as follows:



Figure 2: Plane 0 organization (zones)

TL - Top left rule, sensor node in the top left corner, i.e. origin coordinate (0, 0)
TC - Top center rule, sensor node coordinates (0, 1) to (0, 9)
BC - Bottom center rule, sensor node coordinates (10, 1) to (10, 9)
TR - Top right rule, sensor node coordinate (0, 10)
BL - Bottom left rule, sensor node coordinate (10, 0)
BR - Bottom right rule, sensor node coordinate (10, 10)
Right rule, sensor node coordinates (1, 10) to (9, 10)
Left rule, sensor node coordinates (1, 0) to (9, 0)
The rest of a cell space is WSN rule, i.e. local transition rule of the model.

The initial model was organized as shown in Figure 2. The partitioning of the cell space into 'zones' was done to experiment with the problem at hand (initially) by isolating the 'zones'. Hence, easing the task to be solved, by observing sensors' operations at different 'zones' (and their interaction) prior to optimizing the final solution.

When compared to the initial model, the improved model (refer to Figure 3) are the following:

- Entire cell-space can be populated with sensor nodes
- There are no limitations to functionality of the model (Plane 0 organization)
- Cell space is not divided into zones
- Enables more flexibility to the initial model (reduces significantly the code size)
- Randomness was implemented within the model (it enables us to represent more closely the real-world applications related to WSN)
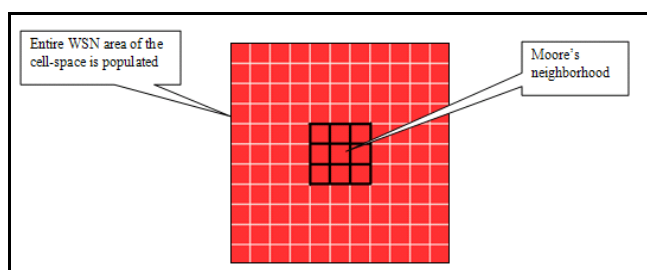


Figure 3: Plane 0 organization (improved model)

In addition to the existing states present in the initial WSN model (described earlier), an additional state was implemented in Plane 0 of the improved WSN model:

**3** – WSN sensor node is typically in ***stand-by mode***; however, nodes randomly become active if one or two nodes are active Within the neighborhood (Moore's neighborhood coverage area); node goes back to ***stand-by mode*** if the condition is true (i.e. one or two neighborhood sensors are active), otherwise it remains in the ***active mode***.

## 2.1 WSN behavior definition

In Plane 0, each Moore's neighborhood (consisting of 9 cells) typically is covered by one active cell (with the value of 2) and the rest of the stand-by cells have a value of 1, whilst the passive cells have a value of 0. At the initial stage all the WSN sensor nodes deployed within the WSN network area (represented by cell space) are active for several time steps until they configure themselves to active and stand-by mode nodes.

In Plane 1, for every time step, the energy of active cells is decreased; the amount of energy available $x$, for any active and stand-by mode cells before cell dies is $0 \leq x \leq 0.8$. Each cells energy level in Plane 1 serves as a reference (memory) of the Plane 0 active and stand-by mode cells (represented by value of 2 and 1 respectively). Plane 0 refers to Plane 1 by (0, 0, 1), while Plane 1 refers to Plane 0 by (0, 0,-1) neighborhood coordinate (refer to Figure 4). The neighboring cells which are in a stand-by mode are represented by the value of 1 within the cell, and are also decreased in negligible quantities during the stand-by mode; The active cells become passive after their energy is consumed and are replaced by the neighboring cells which are yet alive (cells currently in stand-by mode).
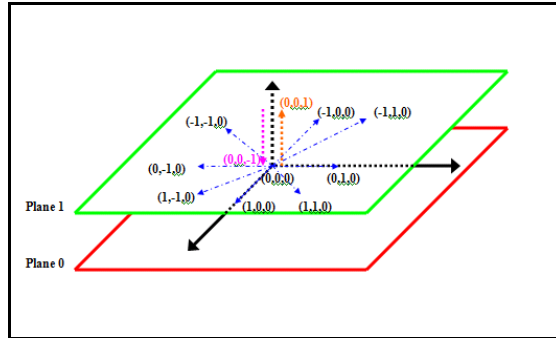
Figure 4: Cell space definition

The neighboring cells which are in a stand-by mode are represented by the value of 1 within the cell, and are also decreased in negligible quantities during the stand-by mode; The active cells become passive after their energy is consumed and are replaced by the neighboring cells which are yet alive (cells currently in stand-by mode). The cell-space in Plane 0 defines sensor nodes, whilst in Plane 1 energy level corresponding to each node (refer to Figure 4).

The algorithm steps through each neighborhood and decides on which nodes stays active and which is to be configured as a stand-by nodes; where each neighborhood of cells is checked for present active cells by comparing the present cell's value with the residual energy (which is check referenced in Plane 1) until they becomes passive (represented by value of 0 in the cell space); Afterwards, one of the stand-by neighboring cell is assigned to becomes active, and this process continues until all the cells become passive (i.e. the energy resources of all sensors is consumed). Plane 1 updates the energy levels of each cell during the simulation as per the specifications.

The WSN simulation model provides closer approximation to the WSN topology algorithm by implementation of randomness within the WSN model; where randomness, redundancy and configuration of nodes play a significant role in reflecting the key factors in deployed networks, such as a network lifetime, coverage area and ratios of active and stand-by sensors at specific points in time. Randomness was implemented by adding another rule (refer to the rule 3 below ) to the model, where stand-by cells randomly become active if only one or two active neighbors are active. The actual results obtained with the improved model, more closely reflect the real-world scenarios and provide better insight into the WSN topology problem.

The problem could be reduced and coded with less than 30 lines of code utilizing the CD++ toolkit and Cell-DEVS, as shown in the following figure

```
[WSN]
type : cell                    dim: (33, 33, 2)                    delay : transport
border : nowrapped
neighbors :  (-1,-1,0)   (-1,0,0)   (-1,1,0)  (0,-1,0)    (0,0,0)    (0,1,0)
neighbors :  (1,-1,0)    (1,0,0)    (1,1,0)   (0,0,1)     (0,0,-1)

localtransition : WSN-rule
[WSN-rule]
rule : { (0,0,0) - 0.0165 }  1000 {cellpos(2) = 1 and  (0,0,-1) = 2 }
rule : { (0,0,0) - 0.00006 } 1000 {cellpos(2) = 1 and (0,0,-1) =1  }
rule : { (0,0,0) - 0.0165 }  1000 {cellpos(2) = 1 and  (0,0,-1) = 3 }
rule : 3 1000 {cellpos(2)=0  and (0,0,1)>0 and
              ( (0,0,0)=1   and (statecount(2)=2 or statecount(2)=1) and randInt(30)=11 ) }
rule : 2 1000 {cellpos(2)=0  and (0,0,0) !=0 and (0,0,0) !=3 and (0,0,1) > 0 and (0,-1,0)!=2
           and (1,-1,0)!=2 and (1,0,0)!=2 and (1,1,0)!=2 }
rule : 1 1000 {cellpos(2)=0  and (0,0,0) !=0 and (0,0,1)>0 and
              ( (0,-1,0)=2  or (1,-1,0)=2 or (1,0,0)=2 or (0,1,0)=2 or (1,1,0)=2 or
              (cellpos(2)=0  and (0,0,0)=3) or (-1,-1,0)=2 or (-1,0,0)=2 or (-1,1,0)=2 ) }
rule : 0 1000 {cellpos(2)=0  and ((0,0,0)=1 or (0,0,0)=2 or (0,0,0)=0)  and (0,0,1) <= 0 }
```

# 3    SIMULATION RESULTS

We executed numerous tests, and in this section we present some of the simulation results obtained and discuss their meaning. The Figure 5 shows the graphical representation of the sensor node states defined in the simulation model:
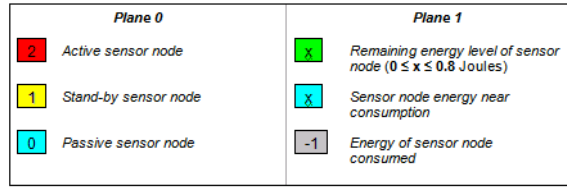
Figure 5: Possible sensor states of the initial WSN model

On the left side of the Figure 5 we have the Plane 0 configuration, where **2** - indicates that a sensor node is ***active***, **1** – indicates that a sensor node is in ***stand-by mode*** and **0** - indicates that a sensor node is in ***passive state*** (i.e. energy of a node is consumed); while on the right side of the Figure 5 we have the Plane 1 configuration, where **x** – is the energy level of WSN sensor nodes (green cells are the sensor nodes with enough energy, while blue color signifies that the sensor nodes are close to dying), **1** - WSN sensor node is ***passive*** (i.e. energy of a node is consumed)

The first simulation results presented are shown in Figure 6 below. Based on the specification, we can see the sensor nodes active at the beginning of the simulation (red cells), while the light blue cells represent the unpopulated zones within the cell space.
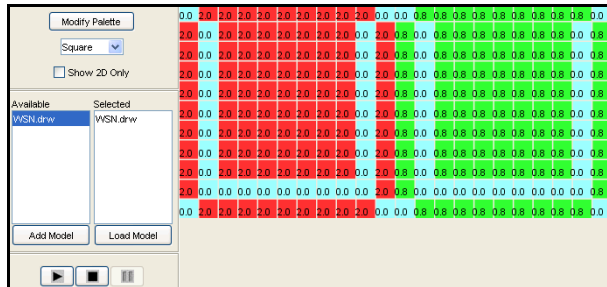


Figure 6: Plane 0 (left, red) and Plane 1 (right, green) prior to execution

The following Figure 7 shows the WSN network during the process of reconfiguration, where nodes are trying to form a structure and every neighbor is trying to set some sensors in active mode while others remain in stand-by. The Plane 1 stores the energy level of each sensor node, which can be monitored.



Figure 7: Snapshot of simulations results after 6 time steps

Figure 8 shows that the WSN sensors are reconfigured as per the specifications and each neighborhood is covered by typically one active node while others remain in stand-by mode (energy of each note is decreasing during each step based on the model specification).

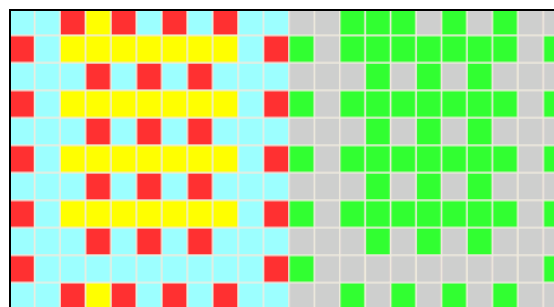Figure 8: Snapshot of simulations results after 15 time units



Figure 9: Snapshot of simulations results after 62 time units

Figure 9 shows that several active nodes (from previous step) are now passive (died cells - Plane 1 on the right, gray cells), while cells that were previously in stand-by mode are taking over by becoming active.

In these simulations, all the sensor nodes become passive after 197 time steps. By comparing the end result of our simulation and the results obtained in (Cunha, Silva, and Loreiro 2005) in particular, the number of active sensors after 200 time units), it can be observed that after approximately 200 time steps all the active sensors become passive (as energy of all nodes is consumed). Hence, the simulation results obtained by our model very closely reflect the same behavior (i.e. after 197 time units the active sensors become passive). These results can be seen in Figure 10.
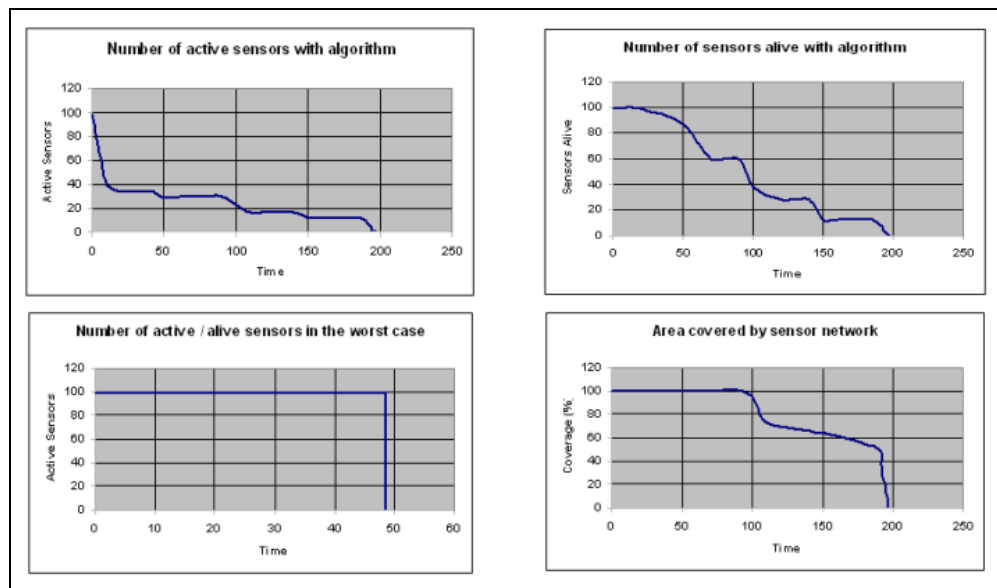


Figure 10: Graphical representation of the Initial WSN simulation results using Cell-DEVS

From the Figure 10 it can be observed that in the worst case scenario (bottom left graph), when all nodes are active all the time, the sensor nodes die after 48.48 time units (where the energy of a node in each step is consumed by 0.0165 J and the initial level of sensors' energy is 0.8 J), while when using the Topology Control Algorithm, where selected sensor nodes within a Moore's neighborhood are alive (top left graph), the number of alive/active sensor nodes decreases gradually, extending the life of sensor nodes to 197 time units (top right graph) and the coverage area (bottom right graph).

Our following results present an improved version of the WSN model, in which we use the following is the graphical representation of the sensor node states defined within the simulation model:
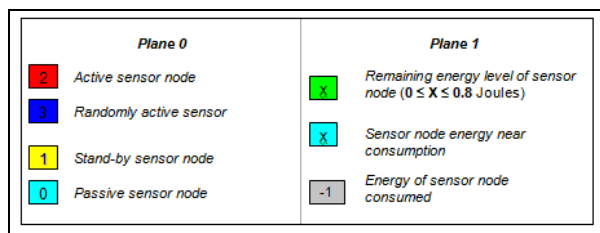
Figure 11: Possible sensor states of the improved WSN model

With respect to the initial model (refer to Figure 5), in representing the sensor node states of the improved model, we have added an addition sensor state to indicate the randomly active sensor nodes, as depicted in Figure 11.

This model provides closer approximation to the WSN topology algorithm by implementation of the randomness within the WSN deployed. The actual results obtained, more closely reflect the real-world scenarios and provide better insight into the WSN topology problem and how efficiently similar problems can be implemented utilizing Cell-DEVS and CD++ tool-kit. The improved model is quite easy to modify in order to simulate different sizes of WSN networks. In this case, it requires to change only the values *n* and *m* (i.e., one line of specification, *dim : (n, m, 2)* where *n* is the number of columns and *m* is the number of rows in the model and provide the desired initial energy levels for the sensor nodes). The example below (refer to Figure 12) represent the simulation model for WSN33, where the cell space is constituted by 33 rows and 33 column, and the total number of cells (sensor nodes) is 1089.
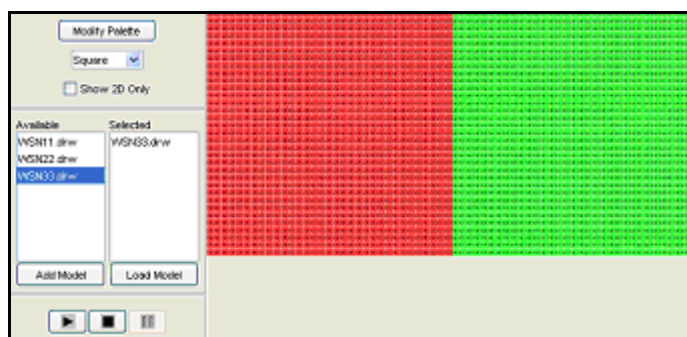


Figure 12: Snapshot of Plane 0 (left, red) and Plane 1 (right, green) prior to execution

The following figure shows the 33 x 33 WSN network during the process of reconfiguration, where nodes are trying to form a structure and every neighbor is trying to set some sensors in active while others remain in stand-by mode. Some of the stand-by nodes are randomly awaken (in Figure 13 left, the cells in blue in Plane 0). In case that only one or two sensor nodes are active, stand-by nodes randomly become active and return to stand-by mode only if one or two more sensor nodes are currently active within the neighborhood. On the Plane 1 data (refer to Figure 13 right), we can see the energy level of each sensor node up to this time step (see in green, indicating that the energy levels x of sensor nodes are: $x \geq 0.05J$).
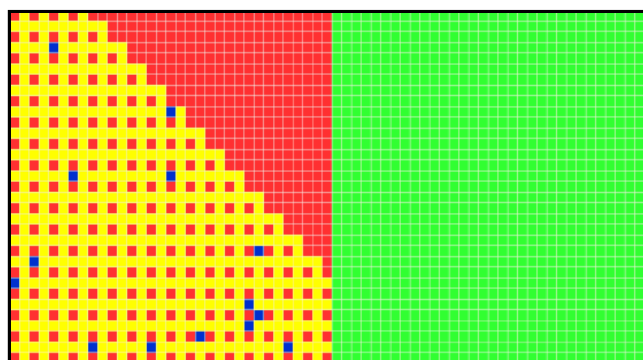


Figure 13: Snapshot of simulations results after 40 time units

Studying the simulation results of this Cell-DEVS model (presented in Figure 14 and Fig 15), we can be see that the coverage area by sensors is reduced after 134 time steps, when more and more nodes become passive (as their energy is consumed). As time progresses, there is a smaller area of the WSN cell-space covered. Finally, after 193 time units, the WSN cell-space becomes passive. Similar, results were obtained when WSN22 and WSN11 were simulated.
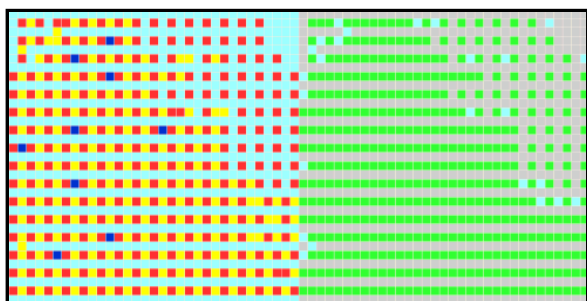


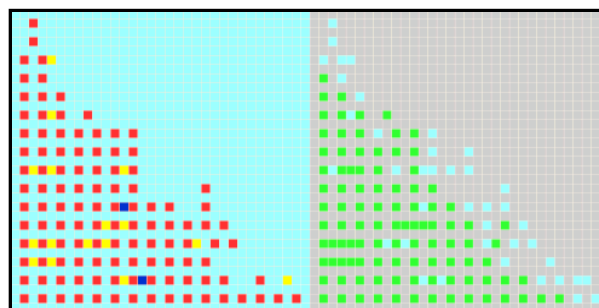Figure 14: Snapshot of simulations results after 134 time units



Figure 15: Snapshot of simulations results after 164 time units

Figure 16 below shows the number of active sensors versus time within the WSN network; evaluated for WSN deployment scenarios within two dimensional cell-spaces, using Cell-DEVS:

a) WSN11 - representing *11x11* cell-space, with 121 sensor nodes
b) WSN22 – representing *22x22* cell-space, with 484 sensor nodes
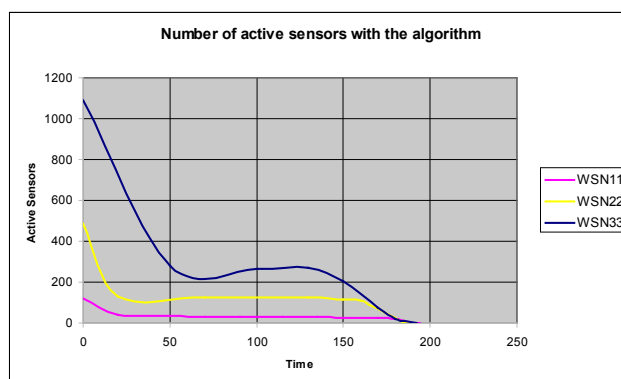c) WSN33 – representing *33x33* cell-space, with 1089 sensor nodes



Figure 16: Number of active sensors with the algorithm using Cell-DEVS

As we can see, the number of active sensors decreases after the configuration of sensor nodes, following the deployment, after which, redundancy is reduced by having only one active node within the Moore's neighborhood (while other nodes are in stand-by). In addition, within each Moore's neighborhood, the stand-by nodes become active randomly and remain active if no sensor is active or return to stand-by mode if any sensor node is still active. The results obtained provide clear indication that network lifetime is increase approximately by 4, which was shown also in (Cunha, Silva, and Loreiro 2005).

On Figure 17, it can be observed that the number of stand-by sensors increases while the WSN cell-space is being configured, and starts to decrease as the active sensors' energy is consumed (refer to Figure 16) hence stand-by sensors become active. Similarly, in Figure 18 we can be observe that the number of sensors alive is at it maximum when the simulation starts (i.e. all the WSN sensor nodes are active) and starts slowly decreasing as the time progresses.
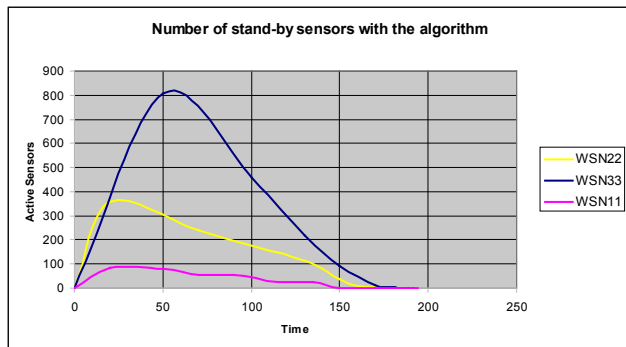
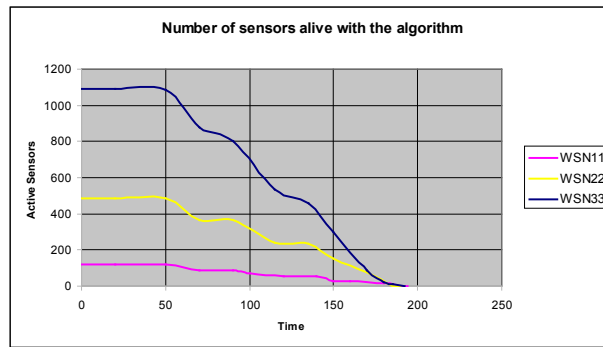Figure 17: Number of stand-by sensors with the algorithm using Cell-DEVS



Figure 18: Number of sensors alive with the algorithm using Cell-DEVS

An important aspect considered in WSN networks is the coverage area, which is closely related to the active sensors within the cell-space. When the redundancy of sensor nodes within the WSN is controlled, the network lifetime is prolonged; hence wider coverage area is maintained for a longer time, as shown by the simulation results in Figure 19.
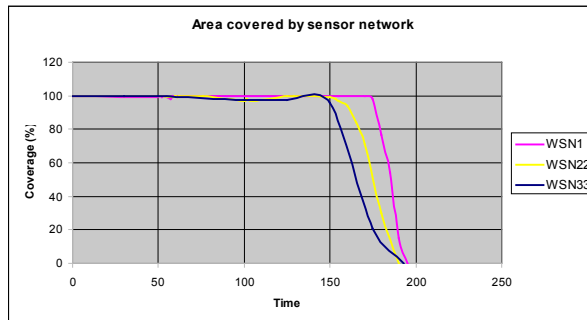


Figure 19: Sensor network coverage area (%) using Cell-DEVS

Finally, in Figure 20 we can see the number of active sensors in the worst case scenario when WSN sensor nodes are active the entire time until the energy of sensors is consumed just before time 50; the network lifetime is approximately 4 time less in comparison to the implementation of WSN topology control algorithm using Cell-DEVS (refer to Figure 16, Figure 18 and Figure 19).
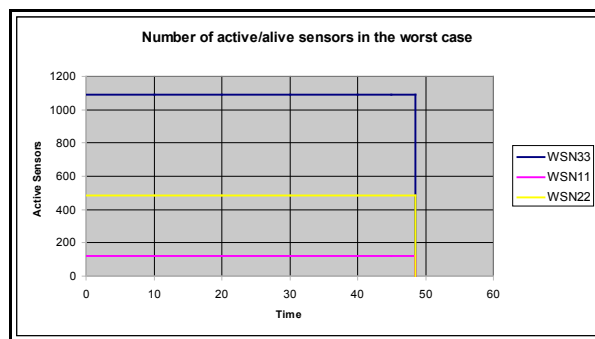


Figure 20: Number of active sensors with the algorithm

## 4 CONCLUSIONS

The objective of this paper was to simulate a large Wireless Sensor Network (WSN) using Cell-DEVS, by implementing the Topology Control Algorithm as presented in. By observing and evaluating the behavior of WSN simulation model, under different test scenarios, it was proven the effectiveness of Cell-DEVS and the CD++ toolkit, as an elegant approach to model, simulate and analyze, in this case the WSN topology problem.

The initial WSN model was further enhanced in order to provide closer approximation to the WSN topology algorithm by implementation of the randomness within the WSN deployed. In addition, modification and simplification of model was done, where cell-space is not divided into zones; allowing more flexibility to model any possible WSN cell-space configuration (entire or partial cell-space populated with sensor node).

The actual results obtained with the improved model, more closely reflect the real-world scenarios and provide a better insight into the WSN topology problem, in particular. As it was observed in section 3, the complexity of the problem can be simplified and coded with less than 30 lines of code, utilizing CD++ toolkit and Cell-DEVS approach; whilst similar problem if implemented in C/C++ (or any other high level programming language) could possibly take up to several hundreds or perhaps thousands of lines of code (based on the implementation approach taken). Thus, proving how efficiently complex problems of similar nature can be implemented simulated and analyzed utilizing CD++ toolkit and Cell-DEVS approach.

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